

Sommes de nombres de Catalan, méandres et polynômes en $1/\pi$

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joint work with Paul Thévenin and Alin Bostan

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Les midis de la Combinatoire

Nancy, 26 novembre 2024



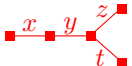
Problem statement

For a tree T , we consider

$$S(T) := \sum_{(x_e) \in \mathbb{Z}_+^{E(T)}} \left(\prod_{v \in V(T)} \text{Cat}_{X_v} 4^{-X_v} \right),$$

where $\text{Cat}_k = \frac{1}{k+1} \binom{2k}{k}$ and $X_v = \sum_{e \ni v} x_e$.

Example

Let $T =$  , then

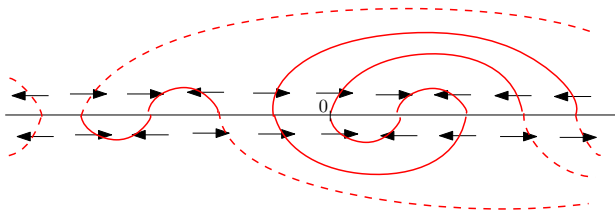
$$S(T) = \sum_{x,y,z,t \geq 0} \text{Cat}_x \text{Cat}_{x+y} \text{Cat}_{y+z+t} \text{Cat}_z \text{Cat}_t 16^{-x-y-z-t}.$$

Our goal: compute these sums, and prove that they are in $\mathbb{Q}[\frac{1}{\pi}]$.

Motivation: meandric systems (1/5)

Definition (The Uniform Infinite Meandric System, or Infinite Noodle)

Draw two bi-infinite sequences of i.i.d. left/right arrows and close them in the unique non-crossing way.



Question

Is there an infinite connected component? What is the distribution of the size of the component of 0? In other words, compute $\mathbb{P}(|C_0| = k)$.

Motivation: meandric system (2/5)

Conjectures

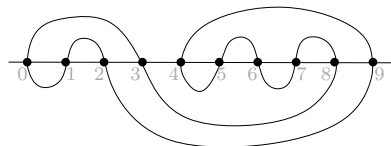
- Almost surely, there is no infinite component (Curien–Kozma–Sidoravicius–Tournier, '19).
- As k tends to $+\infty$, we have $\mathbb{P}(|C_0| \geq k) \sim k^{-(2\sqrt{2}-1)/7+o(1)}$. (Borga–Gwynne–Park, '23).

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A related conjecture: define a **meander** as a **connected finite arc configuration**

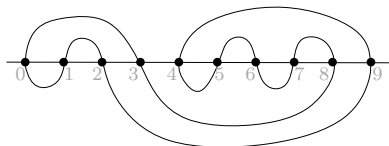


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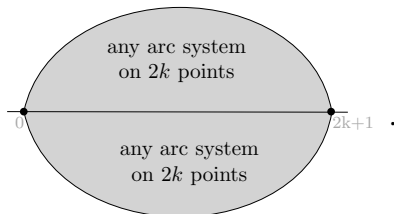
Conjecture (Di Francesco–Golinelli–Guitter, '00):

$$\#\{\text{meanders of size } 2n\} \sim C A^n n^{-\alpha},$$

$$\text{with } \alpha = \frac{29 + \sqrt{145}}{12}.$$

Motivation: meandric system (3/5)

A configuration with $|C_0| = 2$ looks like

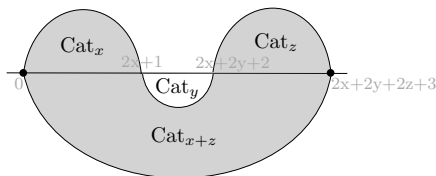


Hence

$$\mathbb{P}[|C_0| = 2] = 2 \sum_{k \geq 1} \text{Cat}_k^2 2^{-4k-4} = \frac{1}{8} S(\blacksquare \dashrightarrow \blacksquare).$$

Motivation: meandric system (4/5)

A configuration with $|C_0| = 4$ looks like



Hence

$$\begin{aligned} \mathbb{P}[|C_0| = 4] &= 4 \times 2 \times \sum_{x,y,z \geq 0} \text{Cat}_x \text{Cat}_y \text{Cat}_z \text{Cat}_{x+z} 2^{-4x-2y-4z-8} \\ &= \frac{1}{32} S(\blacksquare\text{---}\blacksquare\text{---}\blacksquare) S(\square\text{---}). \end{aligned}$$

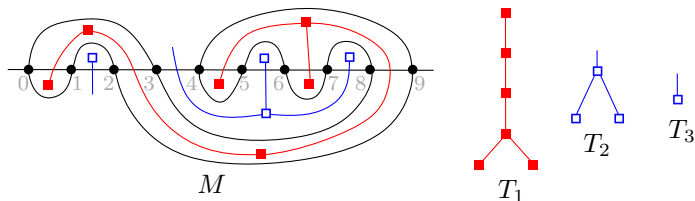
Motivation: meandric system (5/5)

More generally, to compute $\mathbb{P}[|C_0| = k]$,

- we sum over all possible component “shapes”, i.e. over meanders of size k ;
- for a meander M ,

$$\mathbb{P}(S_0 = M) = 2^{-4k+1} k \prod_{i=1}^d S(T_i),$$

where the T_i 's are the “dual trees” of the meander.



Quelques exemples en mathematica

[Link to the notebook](#)

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Results (P_n , P_n^h , S_n are respectively the path, the path with a half-edge and the star with n vertices):

$$S(P_2) = \frac{16}{\pi} - 4;$$

$$S(P_2^h) = \frac{8}{\pi};$$

$$S(P_3) = 8 - \frac{64}{3\pi};$$

$$S(S_4) = \frac{64}{15\pi};$$

$$S(P_4) = -32 + \frac{64}{\pi} + \frac{128}{\pi^2}.$$

$S(P_2)$ and hypergeometric functions

We want to compute $S(P_2) = \sum_{x \in \mathbb{Z}_+} u_x$, where $u_x = (\text{Cat}_x 4^{-x})^2$.

Note: the quotient u_{x+1}/u_x is a rational function in x . Such sums are called hypergeometric. Typical hypergeometric sums are

$${}_2F_1(a, b; c; z) := \sum_{n \geq 0} \frac{a \uparrow^n b \uparrow^n z^n}{c \uparrow^n n!},$$

where $u \uparrow^n := u(u+1) \cdots (u+n-1)$.

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In fact, we have

$$S(P_2) = 4 \cdot {}_2F_1\left(-\frac{1}{2}, -\frac{1}{2}; 1; 1\right) - 4.$$

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Lemma (Gauss identity)

If $c - a - b > 0$, we have

$${}_2F_1(a, b; c; 1) = \frac{\Gamma(c)\Gamma(c-a-b)}{\Gamma(c-a)\Gamma(c-b)}.$$

Thus ${}_2F_1\left(-\frac{1}{2}, -\frac{1}{2}; 1; 1\right) = \frac{4}{\pi}$ and $S(P_2) = \frac{16}{\pi} - 4$.

$S(P_3)$ and the quadratic Catalan recurrence

We want to compute $S(P_3) = \sum_{x,y \in \mathbb{Z}_+} \text{Cat}_x \text{Cat}_y \text{Cat}_{x+y} 16^{-x-y}$.

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Rewrite the sum using $Z = x + y$.

$$\begin{aligned} S(P_3) &= \sum_{Z \geq 0} \text{Cat}_Z 16^{-Z} \left(\sum_{\substack{x,y \geq 0 \\ x+y=Z}} \text{Cat}_x \text{Cat}_y \right) \\ &= \sum_{Z \geq 0} \text{Cat}_Z \text{Cat}_{Z+1} 16^{-Z}. \end{aligned}$$

Looks like $S(P_2)$ with a shift of indices.

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Again, this can be related to hypergeometric functions, namely

$$S(P_3) = 8 - 8 \cdot {}_2F_1\left(-\frac{1}{2}, \frac{1}{2}; 2; 1\right)$$

and Gauss identity allows to compute

$$S(P_3) = 8 - \frac{64}{3\pi}.$$

$S(P_2^h)$ and manipulating inequalities

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By symmetry, we also have

$$S(P_2^h) = \sum_{x \geq z \geq 0} \text{Cat}_x \text{Cat}_z 4^{-x-z}.$$

and thus

$$\begin{aligned} 2S(P_2^h) &= \sum_{x,z \geq 0} \text{Cat}_x \text{Cat}_z 4^{-x-z} + \sum_{\substack{x,z \geq 0 \\ x=z}} \text{Cat}_x \text{Cat}_z 4^{-x-z} \\ &= \left(\sum_{x \geq 0} \text{Cat}_x 4^{-x} \right)^2 + S(P_2) = 4 + \left(\frac{16}{\pi} - 4 \right) = \frac{16}{\pi}. \end{aligned}$$

From edge-variables to vertex variables

We root T and write $v' \leq_T v$ if v' is a descendent of v . Color white (resp. black) vertices at even (resp. odd) distance from the root.

Lemma

The map sending $(x_e)_{e \in E(T)}$ to $(X_v)_{v \in V(T)}$ where $X_v = \sum_{e \ni v} x_e$, is injective

$$(*) \left\{ \begin{array}{l} \sum_{w \in V_o(T)} X_w = \sum_{b \in V_\bullet(T)} X_b \\ \text{for all } v \in V_o(T), \sum_{w \in V_o(T), w \leq_T v} X_w \geq \sum_{b \in V_\bullet(T), b \leq_T v} X_b \\ \text{for all } v \in V_\bullet(T), \sum_{w \in V_o(T), w \leq_T v} X_w \leq \sum_{b \in V_\bullet(T), b \leq_T v} X_b. \end{array} \right.$$

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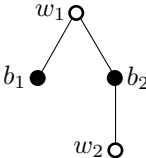
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Corollary

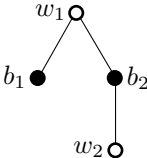
$$S(T) = \sum_{\substack{(X_v) \in \mathbb{Z}_+^{V(T)} \\ (*)}} \prod_{v \in V} \text{Cat}_{X_v} 4^{-V}.$$

Next example

A rooted version of P_4 is . We get

$$S(P_4) = \sum_{\substack{w_1, w_2, b_1, b_2 \geq 0 \\ w_1 + w_2 = b_1 + b_2, \quad b_2 \geq w_2}} \text{Cat}_{w_1} \text{Cat}_{w_2} \text{Cat}_{b_1} \text{Cat}_{b_2} 4^{-w_1 - w_2 - b_1 - b_2}.$$

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Manipulating the inequalities,

$$\begin{aligned} 2S(P_4) &= \sum_{\substack{w_1, w_2, b_1, b_2 \geq 0 \\ w_1 + w_2 = b_1 + b_2}} \text{Cat}_{w_1} \text{Cat}_{w_2} \text{Cat}_{b_1} \text{Cat}_{b_2} 4^{-w_1 - w_2 - b_1 - b_2} \\ &\quad + \sum_{\substack{w_1, w_2, b_1, b_2 \geq 0 \\ w_1 + w_2 = b_1 + b_2, b_2 = w_2}} \text{Cat}_{w_1} \text{Cat}_{w_2} \text{Cat}_{b_1} \text{Cat}_{b_2} 4^{-w_1 - w_2 - b_1 - b_2}. \end{aligned}$$

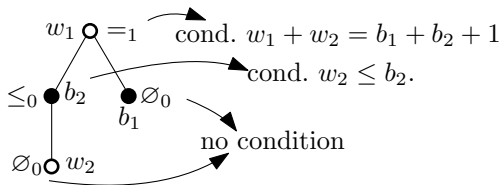
The first term can be computed using the quadratic recurrence, and the second is $S(P_2)^2$.

The general framework

- We consider rooted trees \tilde{T} with black/white/gray* vertices and vertex decorations in $\{=, \leq, \geq, \emptyset\} \times \mathbb{Z}$.
- To each vertex is associated a condition (C_v) comparing the sums of white variables and black variables below it.
- We consider the sum

$$S(\tilde{T}) := \sum_{(X_v) \in \mathbb{Z}_+^{V_{\bullet/o}(\tilde{T})}} \left(\prod_{v \in V_{\bullet/o}(\tilde{T})} \text{Cat}_{X_v} 4^{-X_v} \right) \left(\prod_v 1[C_v] \right).$$

Example: A decorated version of P_4



* Gray vertices do not have associated variables but may impose conditions.

Our main result (general version)

Theorem (Bostan, F., Thévenin, '25)

For any colored decorated tree $S(\tilde{T})$, the sum $S(\tilde{T})$ belongs to $\mathbb{Q}[\frac{1}{\pi}]$.

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For any fixed k , $\mathbb{P}[|C_0| = k]$ is in $\mathbb{Q}[\frac{1}{\pi}]$.

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Corollary

For any fixed k , $\mathbb{P}[|C_0| = k]$ is in $\mathbb{Q}[\frac{1}{\pi}]$.

More generally, we can replace 4^{-X_v} by t^{X_v} and we have

Theorem (Bostan, F., Thévenin, '25)

The series $S(\tilde{T})(t)$ belongs to the space

$$\mathbb{Q}\left[t^{\pm 1}, \sqrt{1-4t}, {}_2F_1\left(-\frac{1}{2}, -\frac{1}{2}; 1; 16t^2\right), {}_2F_1\left(-\frac{1}{2}, \frac{1}{2}; 2; 16t^2\right)\right].$$

Moreover, ${}_2F_1\left(-\frac{1}{2}, -\frac{1}{2}; 1; 16t^2\right)$ and ${}_2F_1\left(-\frac{1}{2}, \frac{1}{2}; 2; 16t^2\right)$ are algebraically independent over $\mathbb{Q}[t^{\pm 1}, \sqrt{1-4t}]$.

Some words on the proof

We have some reduction rules, like

$$\begin{array}{c}
 \begin{array}{cccc}
 \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} (\geq, K) & + & \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} (\leq, K) & \equiv & \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} (=, K) & + & \begin{array}{c} \diagup \\ \bullet \\ \diagdown \end{array} (\emptyset, K) \\
 \end{array}
 \end{array}
 ,$$

(manipulating inequalities)

or

$$\begin{array}{c}
 \begin{array}{c} \triangle \\ U \end{array} \begin{array}{c} \Delta \\ \circ \\ \circ \\ \emptyset_0 \\ \triangle \\ V \end{array} \equiv t^{-1} \cdot \begin{array}{c} \triangle \\ U \end{array} \begin{array}{c} \Delta+1 \\ \circ \\ \triangle \\ V \end{array} - t^{-1} \cdot \begin{array}{c} \triangle \\ U \end{array} \begin{array}{c} \Delta+1 \\ \bullet \\ \triangle \\ V \end{array} .
 \end{array}$$

(using the quadratic recurrence)

Reducing any tree with such rules?

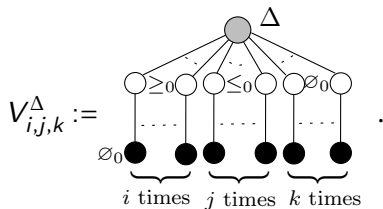
Let T be a bicolored tree. We start from the leaves.

- If there are two twin leaves of the same color or if one leaf has the same color than its parents, we can apply the quadratic recurrence 😊

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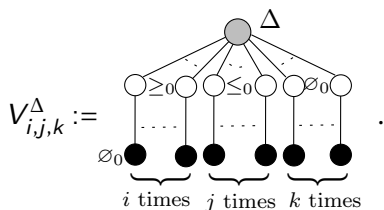
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- The difficult case is that of **long stars**, e.g.



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By *reversing inequalities*, we obtain relations between the $V_{i,j,k}^{\Delta}$ for various values of i , j and k .

A linear system for long stars

Lemma

For any “rootstock” R , any decoration Δ and any $d \geq 2$, we have

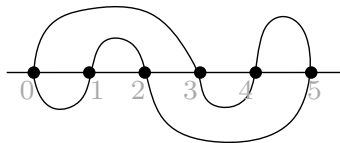
$$\begin{pmatrix} 2 & 1 & 0 & \dots & 0 \\ 1 & 2 & 1 & \ddots & \vdots \\ 0 & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & 1 & 2 & 1 \\ 0 & \dots & 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} S(R | V_{1,d-1,0}^\Delta) \\ \vdots \\ S(R | V_{d-1,1,0}^\Delta) \end{pmatrix} = \begin{pmatrix} X_1 \\ \vdots \\ X_{d-1} \end{pmatrix} - \begin{pmatrix} S(R | V_{0,d,0}^\Delta) \\ 0 \\ \vdots \\ 0 \\ S(R | V_{d,0,0}^\Delta) \end{pmatrix},$$

where, for $1 \leq i \leq d-1$,

$$X_i = S(R | V_{i-1,d-1-i,2}^\Delta) + 2S(R | V_{i-1,d-1-i,1}^\Delta) \cdot S_{=,0} + S(R | V_{i-1,d-1-i,0}^\Delta) \cdot S_{=,0}^2.$$

→ this lemma allows to express $S(R | V_{i,j,0}^\Delta)$ in terms of smaller trees since the matrix is invertible!

Thanks for your attention!



$$\sum_{x,y,z} \text{Cat}_x \text{Cat}_{x+y} \text{Cat}_{y+z} \text{Cat}_z 16^{-x-y-z} = -32 + \frac{64}{\pi} + \frac{128}{\pi^2}.$$

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