The WZ phenomenon (Part 1)

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The WZ method: The Algorithm

- 1. Suppose we want to show $\sum_{k} t(n, k) = RHS(n)$.
- 2. If RHS(n) $\not\equiv$ 0, look instead at $f(n) := \sum_k F(n, k) = 1$, where $F(n, k) = \frac{t(n, k)}{\mathsf{RHS}(n)}$.
- 3. Find R(n, k), the certificate for the WZ method (see later, steps 3.1 and 3.2) and define G(n, k) := R(n, k)F(n, k).
- 4. Check that D:=F(n+1,k)-F(n,k)=G(n,k+1)-G(n,k) (mostly easily verifiable). Summing this equation over all k and noting that the RHS telescopes to 0, we get $\sum_k F(n+1,k) = \sum_k F(n,k), \text{ i.e. } \sum_k F(n,k) \text{ doesn't depend on } n \text{ and is therefore constant.}$
- 5. Verify that this constant is indeed 1 by entering an explicit value for n, e.g. check that $\sum_{k} F(0, k) = 1$.



The WZ method: WZ pair

Definition

A pair (F(n, k), G(n, k)) is called a **WZ pair**, if it fulfills

$$F(n+1,k) - F(n,k) = G(n,k+1) - G(n,k)$$

The WZ method: Step 3: Find R(n, k)

For the crucial Step 3 of the WZ method we use Gosper's algorithm in the following way:

- 3.1 Let D(k) := F(n+1,k) F(n,k) and input D(k) into Gosper's algorithm. If Gosper fails, then the WZ method fails as well.
- 3.2 Otherwise Gosper gives us a function g(k) such that D(k) = g(k+1) g(k). Of course this g will also contain a parameter n, so rename it to G(n,k). This G(n,k) is the WZ mate for F(n,k). Furthermore $\frac{G(n,k)}{F(n,k)} := R(n,k)$ is a rational function.

The WZ method: Do we have a problem?

As seen in a previous presentation: for every proper hypergeometric term F(n,k) we always have a telescoping certification of $\sum_k F(n,k) = \text{const.}$ as

$$\sum_{j=0}^{J} a_j(n) F(n+j,k) = G(n,k+1) - G(n,k)$$

If we divide by the RHS (if RHS \neq 0), then very often this LHS just reduces to F(n+1,k) - F(n,k).

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The WZ method is valid: Theorem

Theorem

Let (F, G) be a WZ pair, i.e. s.t.

$$F(n + 1, k) - F(n, k) = G(n, k + 1) - G(n, k)$$
. Assume

(G1)
$$\forall n \in \mathbb{N} : \lim_{k \to \pm \infty} G(n, k) = 0$$

Then $\sum_{k} F(n, k) = const. \ \forall n \in \mathbb{N}$, i.e. the certification procedure is valid.

The WZ method is valid: Proof

Let $\Delta_n h(n) := h(n+1) - h(n)$. Starting from the WZ equation and summing over $-L \leqslant k \leqslant K$ we get:

$$\sum_{k=-L}^{K} (F(n+1,k) - F(n,k)) = \sum_{k=-L}^{K} (G(n,k+1) - G(n,k))$$

But the summand on the LHS is just $\Delta_n(F(n,k))$ and the RHS telescopes to G(n,K+1)-G(n,-L). Therefore we get:

$$\Delta_n\left(\sum_{k=-L}^K F(n,k)\right) = G(n,K+1) - G(n,-L)$$

The WZ method is valid: Proof

Now, taking the limits as $K, L \to +\infty$ we get on the LHS $\Delta_n(\sum_k F(n,k)) = \sum_k F(n+1,k) - \sum_k F(n,k)$ and on the RHS, by (G1), we get 0. So overall:

$$\sum_{k} F(n+1,k) = \sum_{k} F(n,k)$$

This means, that $\sum_{k} F(n, k)$ is independent of n, i.e. it is constant.

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Mathematica session: The WZ module

```
(*WZ::usage=
 "WZ[f,n,k] yields the WZ certificate of f[n,k]. Here the input f is an
 expression, not a function. If verbose=TRUE, R(n,k) and G(n,k) are printed.
   If check=TRUE, 3 checks are done and printed. For that the bounds of the
   supports of F (sF, eF) and G (sG, eG) are needed.
 R(n,k) and G(n,k) are returned as expressions."*)
WZnew[f , n , k , verbose , check , sG , eG , sF , eF ] :=
  Module[{df, r, g, WZcheck, TelCheck, ConstCheck},
   df = (f /, \{n \to n + 1\}) - f; (*D(k) = F(n+1,k) - F(n,k) *)
   (*Input D(k) into Gosper's algorithm --> Output: g(k) s.t. g(k+1)-g(k)=D(k)*)
   g = FactorialSimplifv(GosperSum(df, k11:
   r = FactorialSimplify[g/f]; (*R(n,k)=G(n,k)/F(n,k)*)
   If [verbose == TRUE,
    Print("The rational function R(n.k) is ", r1:
    Print["The WZ mate G(n,k) is ", g];
   1;
   If [check == TRUE.
    (*Need F(n+1,k)-F(n,k)=G(n,k+1)-G(n,k) for valid WZ pair*)
    WZcheck = FactorialSimplify[(f /. \{n \rightarrow n+1\}) - f - (g /. \{k \rightarrow k+1\}) + g];
    Print("Check that (F.G) is a valid WZ pair (should give 0): ", WZcheck1:
    (*Need RHS of WZ equation to telescope to 0. Then Sum(F(n,k)) is independent
     of n. i.e. constant*)
    TelCheck = Sum\{(g /, \{k \rightarrow k+1\}) - g, \{k, sG, eG\}\}\}:
    Print["Check that RHS telescopes to 0 (should give 0): ", TelCheck];
    (*Check that this constant is 1 by entering some value for n, e.g. n=1*)
    ConstCheck = Sum[(f /, \{n \rightarrow 1\}), \{k, sF, eF\}];
    Print["Check that Sum(F(n,k))=1 (should give 1):", ConstCheck];
   1;
   \{r, g\} (*Return R(n,k) and G(n,k)*)
                                                                                      1;
```

For $b \in \mathbb{Z}_- \cup \{0\}$ or Re(c-a-b) > 0 we have:

$$_{2}F_{1}\begin{bmatrix} a, b \\ c \end{bmatrix} = \frac{\Gamma(c-a-b)\Gamma(c)}{\Gamma(c-a)\Gamma(c-b)}$$

By definition we can rewrite the LHS as:

$${}_{2}F_{1}\begin{bmatrix} a,b\\c \end{bmatrix} = \sum_{k=0}^{+\infty} \frac{(a)_{k}(b)_{k}}{(c)_{k}} \frac{1^{k}}{k!}$$
$$= \sum_{k=0}^{+\infty} \frac{(a+k-1)!(b+k-1)!(c-1)!}{(a-1)!(b-1)!(c+k-1)!k!}$$

For $b \in \mathbb{Z}_- \cup \{0\}$ or Re(c-a-b) > 0 we have:

$$_{2}F_{1}\begin{bmatrix} a, b \\ c \end{bmatrix} = \frac{\Gamma(c-a-b)\Gamma(c)}{\Gamma(c-a)\Gamma(c-b)}$$

Now, the LHS we can rewrite as:

$$\sum_{k=0}^{+\infty} \frac{(a+k-1)!(b+k-1)!(c-1)!}{(a-1)!(b-1)!(c+k-1)!k!}$$

Then, dividing this new LHS by the RHS and setting a = n to get an identity of the form $\sum_{k} F(n, k) = 1$ we get:

$$\sum_{k=0}^{+\infty} \frac{(n+k-1)!(b+k-1)!(c-1)!\Gamma(c-n)\Gamma(c-b)}{(n-1)!(b-1)!(c+k-1)!k!\Gamma(c-n-b)\Gamma(c)}$$

$$= \sum_{k=-1}^{+\infty} \frac{(n+k)!(b+k)!(c-1)!\Gamma(c-n)\Gamma(c-b)}{(n-1)!(b-1)!(c+k)!(k+1)!\Gamma(c-n-b)\Gamma(c)} \stackrel{!}{=} 1$$

$$\sum_{k=-1}^{+\infty} \frac{(n+k)!(b+k)!(c-1)!\Gamma(c-n)\Gamma(c-b)}{(n-1)!(b-1)!(c+k)!(k+1)!\Gamma(c-n-b)\Gamma(c)} = 1$$

This summand we can now enter in Mathematica:

```
 f := ((n+k)!(b+k)!(c-1)!Gamma[c-n]Gamma[c-b])/ \\ ((n-1)!(b-1)!(c+k)!(k+1)!Gamma[c-n-b]Gamma[c]); \\ WZnew[f, n, k, TRUE, FALSE, , , ,];
```

And we get out:

```
The rational function R(n,k) is -\frac{(1+k)(c+k)}{(-1+c-n)n}
The WZ mate G(n,k) is -\frac{(-1-b+c)!(b+k)!(-2+c-n)!(k+n)!}{(-1+b)!k!(-1+c+k)!(-1-b+c-n)!n!}
```

We want to check that things worked correctly.

```
WZnew[f, n, k, FALSE, TRUE, -1, Infinity, -1,
Infinity];
And indeed we get out:
Check that (F,G) is a valid WZ pair (should give
0): 0
Check that RHS telescopes to 0 (should give 0):
0
Check that Sum(F(n,k))=1 (should give 1):1
```

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The companion identity: Theorem

Theorem

Let (F, G) be a WZ pair, i.e. s.t.

$$F(n+1,k) - F(n,k) = G(n,k+1) - G(n,k)$$
. Assume

(F1) $\forall k$ in the support of F(n, k): $f_k := \lim_{n \to +\infty} F(n, k)$ exists and is finite

(G2)
$$\lim_{L\to +\infty} \sum_{n=0}^{+\infty} G(n, -L) = 0$$

Then $\sum_{n=0}^{+\infty} G(n,k) = \sum_{j=-\infty}^{k-1} (f_j - F(0,j))$, which we call the companion identity.

The companion identity: Proof

Let $\Delta_n h(n) := h(n+1) - h(n)$. Starting from the WZ equation and summing over $0 \le n \le N$ we get:

$$\sum_{n=0}^{N} (F(n+1,k) - F(n,k)) = \sum_{n=0}^{N} (G(n,k+1) - G(n,k))$$

But the LHS telscopes to F(N+1,k)-F(0,k) and the summand on the RHS is just $\Delta_k(G(n,k))$. Therefore we get:

$$F(N+1,k) - F(0,k) = \Delta_k \left(\sum_{n=0}^N G(n,k) \right)$$

The companion identity: Proof

Now, taking the limits as $N \to +\infty$ we get on the LHS, by (F1) $f_k - F(0,k)$ and on the RHS $\Delta_k(\sum_{n=0}^{+\infty} G(n,k))$. If we then sum over $-L \leqslant j \leqslant k-1$ we get overall:

$$\sum_{j=-L}^{k-1} (f_j - F(0,j)) = \sum_{j=-L}^{k-1} \Delta_j \left(\sum_{n=0}^{+\infty} G(n,j) \right)$$

But switching the summation signs on the RHS we get $\sum_{n=0}^{+\infty} \sum_{j=-L}^{k-1} (G(n,j+1) - G(n,j)) = \\ \sum_{n=0}^{+\infty} (G(n,k) - G(n,-L)).$ Then taking the limits as $L \to +\infty$ we get on the LHS $\sum_{j=-\infty}^{k-1} (f_j - F(0,j))$ and on the RHS, by (G2), we get $\sum_{n=0}^{+\infty} G(n,k)$. So overall:

$$\sum_{j=-\infty}^{k-1} (f_j - F(0,j)) = \sum_{n=0}^{+\infty} G(n,k)$$

The companion identity: General version

Adapting the proof slightly (summing over $I \leq n \leq N$ in the beginning, instead of $0 \leq n \leq N$), we also have the more general companion identity

$$\sum_{n=1}^{+\infty} G(n,k) = \sum_{j=-\infty}^{k-1} (f_j - F(I,j))$$

Consider the identity $\sum_k \binom{n}{k}^2 = \binom{2n}{n}$. By dividing through the RHS we get $F(n,k) = \binom{n}{k}^2/\binom{2n}{n}$.

First, we enter this in Mathematica again to find R(n, k) and G(n, k):

$$R(n,k) = \frac{-3n+2k-1}{2(2n+1)}, \ G(n,k) = \frac{(-3n+2k-1)\binom{n}{k}^2}{2(2n+1)\binom{2n}{n}}$$

Using Stirling's formula we get (F1) for $F(n,k) = \binom{n}{k}^2 / \binom{2n}{n}$:

$$f_k := \lim_{n \to +\infty} F(n, k) = \dots = 0$$

$$(G2): \lim_{L \to +\infty} \sum_{n=0}^{+\infty} G(n, -L) = \lim_{L \to +\infty} \sum_{n=0}^{+\infty} \frac{-3n - 2L - 1}{2(2n+1)} \frac{\binom{n}{-L}^2}{\binom{2n}{n}} = 0$$

Note that
$$F(0, k) = \begin{cases} 1, & \text{if } k = 0 \\ 0, & \text{else} \end{cases}$$

Therefore we get as companion identity

$$\left(\sum_{n=0}^{+\infty} G(n,k) = \sum_{i=-\infty}^{k-1} (f_i - F(0,j))\right):$$

$$\sum_{n=0}^{+\infty} \frac{-3n+2k-1}{2(2n+1)} \frac{\binom{n}{k}^2}{\binom{2n}{n}} = -1 \iff \sum_{n=0}^{+\infty} \frac{3n-2k+1}{2n+1} \frac{\binom{n}{k}^2}{\binom{2n}{n}} = 2$$

$$\sum_{n=0}^{+\infty} \frac{3n - 2k + 1}{2(2n+1)} \frac{\binom{n}{k}^2}{\binom{2n}{n}} = 1$$

We can in fact check this new identity (note that we switch the roles of n and k).

```
f := ((3k-2n+1) Binomial[k, n]^ 2)/(2(2k+1)
Binomial[2k, k]);
WZnew[f, n, k, TRUE, FALSE, , , ,];
```

And we get:

The rational function R(n,k) is
$$-\frac{2(1+2k)(k-n)^2}{(1+3k-2n)(1+n)^2}$$

The WZ mate G(n,k) is $-\frac{(k-n)^2}{(1+n)^2} = \frac{2(1+2k)(k-n)^2}{(1+3k-2n)(1+n)^2}$

```
Then to check it:
WZnew[f, n, k, FALSE, TRUE, n, Infinity, 1,
Infinity];
And after some computation time we are rewarded with:
Check that (F,G) is a valid WZ pair (should give
0): 0
Check that RHS telescopes to 0 (should give 0):
0
Check that Sum(F(n,k))=1 (should give 1):1
```

Consider the following identity:

$$\sum_{k=1}^{+\infty} \frac{(n-i)!(n-j)!(i-1)!(j-1)!}{(n-1)!(k-1)!(n-i-j+k)!(i-k)!(j-k)!} = 1$$

To have well-defined factorials, we get the constraints:

$$n \geqslant \max\{i, j, 1\}, \ 1 \leqslant k \leqslant \min\{i, j\}, \ n + k \geqslant i + j$$
 (1)

Enter F(n, k) in Mathematica to find R(n, k) and G(n, k): f := ((n-i)!(n-j)!(i-1)!(j-1)!)/ ((n-1)!(k-1)!(n-i-j+k)!(i-k)!(j-k)!);WZnew[f, n, k, TRUE, FALSE, , , ,];

And we get out:

The rational function R(n,k) is $\frac{-1+k}{n}$ The WZ mate G(n,k) is $\frac{(-1+i)!(-1+j)!(-i+n)!(-j+n)!}{(i-k)!(j-k)!(-2+k)!n!(-i-j+k+n)!}$

```
Then to check it:
WZnew[f, n, k, FALSE, TRUE, 1, Infinity, 1,
Infinity];
And we are indeed rewarded with:
Check that (F,G) is a valid WZ pair (should give
0): 0
Check that RHS telescopes to 0 (should give 0):
0
Check that Sum(F(n,k))=1 (should give 1):1
```

Using Stirling's formula we get (F1) for

$$F(n,k) = \frac{(n-i)!(n-j)!(i-1)!(j-1)!}{(n-1)!(k-1)!(n-i-j+k)!(i-k)!(j-k)!}$$
as $f_k := \lim_{n \to +\infty} F(n,k) = \dots = \begin{cases} 1, & \text{if } k = 1 \\ 0, & \text{if } k \ge 2 \end{cases}$

$$(G2): \lim_{L \to +\infty} \sum_{n=0}^{+\infty} G(n,-L)$$

$$= \lim_{L \to +\infty} \frac{(i-1)!(j-1)!}{(-L-2)!(i+L)!(j+L)!} \sum_{n=0}^{+\infty} \frac{(n-i)!(n-j)!}{n!(n-i-j-L)!} = 0$$

$$F(j,k) = \frac{(j-i)!(i-1)!}{(k-1)!(k-i)!(i-k)!(j-k)!} = \binom{j-i}{j-k} \binom{i-1}{i-k}$$

From constraint $n + k \ge i + j$, for n = j, we get $k \ge i$. From second binomial coefficient we need $k \le i$. So in total k = i.

Therefore we get as companion identity (for l = j):

$$\left(\sum_{n=j}^{+\infty}G(n,k)=\sum_{k'=-\infty}^{k-1}(f_{k'}-F(j,k'))\right)$$

$$\tfrac{(i-1)!(j-1)!}{(k-2)!(i-k)!(j-k)!} \sum_{n=j}^{+\infty} \tfrac{(n-i)!(n-j)!}{n!(n-i-j+k)!} = 1 - \sum_{k'=-\infty}^{k-1} \binom{j-i}{j-k'} \binom{i-1}{i-k'}$$



$$\tfrac{(i-1)!(j-1)!}{(k-2)!(i-k)!(j-k)!} \sum_{n=j}^{+\infty} \tfrac{(n-i)!(n-j)!}{n!(n-i-j+k)!} = 1 - \sum_{k'=-\infty}^{k-1} \binom{j-i}{j-k'} \binom{i-1}{i-k'}$$

For RHS: need $0 \le j - k' \le j - i$ and $0 \le i - k' \le i - 1$, i.e. $j \ge k' \ge i$ and $i \ge k' \ge 1$, i.e. k' = i. But the sum only goes up to k - 1 = i - 1. So the sum in the RHS is 0. Overall:

$$\sum_{n=j}^{+\infty} \frac{(n-i)!(n-j)!}{n!(n-i-j+k)!} = \frac{(k-2)!(i-k)!(j-k)!}{(i-1)!(j-1)!}$$

If we now write n-i-j+k=r in the sum on the LHS and set $i=c-a,\ j=c-b$ and k=c-a-b+1, we get

$$\sum_{r=0}^{+\infty} \frac{(r+a-1)!(r+b-1)!}{r!(r+c-1)!} = \frac{(c-a-b-1)!(b-1)!(a-1)!}{(c-a-1)!(c-b-1)!}$$

which is exactly Gauss's $_2F_1$ identity.

